

Analysis and Design of Algorithms Lecture 5

Sorting Algorithms II

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Counting Sort

Radix Sort

Merge Sort

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Counting Sort

Counting sort is a sorting technique based on keys between a specific range. It works by counting the number of objects having distinct key values (kind of hashing). Then doing some arithmetic to calculate the position of each object in the output sequence.

Counting Sort

Algorithm:

- Step1: Create a count array to store the count of each unique object
- Step2 : Modify count array by adding the previous number.
- Step3 : Create output array by decrease count array

Counting Sort

□ Example 1 Assume the following Array in range of 0 to 5:

1	4	3	2	3	5	2

	1	4	3	2	3	5	2
--	---	---	---	---	---	---	---

|--|

0	1	2	3	4	5
0	0	0	0	0	0



|--|



|--|







1	4	3	2	3	5	2



1	4	3	2	3	5	2





|--|



1 4 3	2	3	5	2
-------	---	---	---	---





1	4 3	2	3	5	2
---	-----	---	---	---	---



1	4 3	2	3	5	2
---	-----	---	---	---	---





1	4	3	2	3	5	2





|--|



|--|





1	4 3	2	3	5	2
---	-----	---	---	---	---



1	4 3	2	3	5	2
---	-----	---	---	---	---



1 4	3	2	3	5	2
-----	---	---	---	---	---

0 1 2 3 4 5 0 1 2 2 1 1

|--|



1	4	3	2	3	5	2



1 4 3 2 3 5	2
-------------	---



1 4 3	2	3	5	2
-------	---	---	---	---



|--|



Counting Sort

Output each object from the input sequence followed by decreasing its count by 1:

1	4	3	2	3	5	2

0	1	2	3	4	5
0	1	3	5	6	7
















Output each object from the input sequence followed by decreasing its count by 1:



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Array is now sorted



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D Python

Code

```
def countSort(arr):
    outputarr = [0 for i in range(127)]
    countarr = [0 for i in range(127)]
```

```
for i in arr:
    countarr[ord(i)] += 1
```

for i in range(127):
 countarr[i] += countarr[i-1]

for i in range(len(arr)):
 outputarr[countarr[ord(arr[i])]-1] = arr[i]
 countarr[ord(arr[i])] -= 1

return outputarr[0:len(arr)]

ASCII Table

Dec	Hex	0ct	Char	Dec	Hex	0ct	Char	Dec	Hex	0ct	Char	Dec	Hex	0ct	Char
0	0	0		32	20	40	[space]	64	40	100	@	96	60	140	`
1	1	1		33	21	41	!	65	41	101	Ā	97	61	141	а
2	2	2		34	22	42		66	42	102	В	98	62	142	b
3	3	3		35	23	43	#	67	43	103	С	99	63	143	с
4	4	4		36	24	44	\$	68	44	104	D	100	64	144	d
5	5	5		37	25	45	%	69	45	105	E	101	65	145	е
6	6	6		38	26	46	&	70	46	106	F	102	66	146	f
7	7	7		39	27	47		71	47	107	G	103	67	147	g
8	8	10		40	28	50	(72	48	110	Н	104	68	150	h
9	9	11		41	29	51)	73	49	111	I	105	69	151	i
10	А	12		42	2A	52	*	74	4A	112	J	106	6A	152	j
11	В	13		43	2B	53	+	75	4B	113	К	107	6B	153	k
12	С	14		44	2C	54	,	76	4C	114	L	108	6C	154	I
13	D	15		45	2D	55	-	77	4D	115	М	109	6D	155	m
14	E	16		46	2E	56		78	4E	116	N	110	6E	156	n
15	F	17		47	2F	57	/	79	4F	117	0	111	6F	157	0
16	10	20		48	30	60	0	80	50	120	Р	112	70	160	р
17	11	21		49	31	61	1	81	51	121	Q	113	71	161	q
18	12	22		50	32	62	2	82	52	122	R	114	72	162	r
19	13	23		51	33	63	3	83	53	123	S	115	73	163	S
20	14	24		52	34	64	4	84	54	124	Т	116	74	164	t
21	15	25		53	35	65	5	85	55	125	U	117	75	165	u
22	16	26		54	36	66	6	86	56	126	V	118	76	166	v
23	17	27		55	37	67	7	87	57	127	W	119	77	167	w
24	18	30		56	38	70	8	88	58	130	Х	120	78	170	х
25	19	31		57	39	71	9	89	59	131	Y	121	79	171	У
26	1A	32		58	3A	72	:	90	5A	132	Z	122	7A	172	Z
27	1B	33		59	3B	73	;	91	5B	133	[123	7B	173	{
28	1C	34		60	3C	74	<	92	5C	134	١	124	7C	174	
29	1D	35		61	3D	75	=	93	5D	135]	125	7D	175	}
30	1E	36		62	3E	76	>	94	5E	136	^	126	7E	176	~
31	1F	37		63	3F	77	?	95	5F	137	_	127	7F	177	

```
arr = "mynameiskhan"
ans = countSort(arr)
print("".join(ans))
```

Time Complexity: O(n+k) where n is the number of elements in input array, and k is the range of input.

D Example of worst case

Range between 1 to 10K

10	5	10k	5k	200	
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Radix Sort

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Radix sort is an algorithm that sorts numbers by processing digits of each number either starting from the least significant digit (LSD) or starting from the most significant digit (MSD).

□ The idea of Radix Sort is to do digit by digit sort starting from least significant digit to most significant digit. Radix sort uses counting sort as a subroutine to sort.
Algorithm:

- Step1: Take the least significant digit of each element
- Step2 : Sort the list of elements based on that digit
- Step3 : Repeat the sort with each more significant digit

Assume the following Array:



□ The Sorted list will appear after three steps

170	45	75	90	802	24	2	66
-----	----	----	----	-----	----	---	----

170	90	802	2	24	45	75	66
-----	----	-----	---	----	----	----	----

802	2	24	45	66	170	75	90

2	24	45	66	75	90	170	802
---	----	----	----	----	----	-----	-----

□ Step1: Sorting by least significant digit (1s place)



170	90	802	2	24	45	75	66

□ Step2: Sorting by next digit (10s place)

1 <u>7</u> 0	<u>9</u> 0	8 <mark>0</mark> 2	2	<u>2</u> 4	<u>4</u> 5	<u>7</u> 5	<u>6</u> 6
--------------	------------	--------------------	---	------------	------------	------------	------------

802	2	24	45	66	170	75	90
-----	---	----	----	----	-----	----	----

□ Step3: Sorting by most significant digit (100s place)

<u>8</u> 02	2	24	45	66	<u>1</u> 70	75	90
-------------	---	----	----	----	-------------	----	----

2	24	45	66	75	90	170	802
---	----	----	----	----	----	-----	-----

Array is now sorted

2	24	45	66	75	90	170	802
---	----	----	----	----	----	-----	-----

Example 2

1	2	3		1	2	3		1	2	3		1	2	3
5	8	3		5	8	3		6	2	5		1	5	4
1	5	4		1	5	4		1	5	4		4	5	6
5	6	7	1 s	6	2	5	10 s	4	5	6	100s	5	6	7
6	8	9		4	5	6		5	6	7		5	8	3
6	2	5		5	6	7		5	8	3		6	2	5
4	5	6		6	8	9		6	8	9		6	8	9

```
def countingSort(arr, count1):
    n = len(arr)
    output = [0] * (n)
    count = [0] * (10)
    for i in range(0, n):
        index = (arr[i]/count1)
        count[ int((index)%10) ] += 1
    for i in range(1,10):
        count[i] += count[i-1]
    i = n - 1
    while i>=0:
        index = (arr[i]/count1)
        output[ count[ int((index)%10) ] - 1] = arr[i]
        count[ int((index)%10) ] -= 1
        i -= 1
    return output
```

□ Python Code

```
def radixSort(arr):
    # Find the maximum number to know number of digits
    maxnum = max(arr)
    count = 1
    while maxnum/count > 0:
        arr=countingSort(arr,count)
        count *= 10
    return arr
```

□ Python Code

arr = [170, 45, 75, 90, 802, 24, 2, 66] print(radixSort(arr))

Time Complexity: O(n+k/d) where n is the number of elements in input array, k is the range of input, and d is number of digits.

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Merge Sort is a Divide and Conquer algorithm. It divides input array in two halves, calls itself for the two halves and then merges the two sorted halves.

Algorithm:

- Step1: Divide the list recursively into two halves until it can no more be divided
- Step2 : Merge (Conquer) the smaller lists into new list in sorted order

Assume the following Array:

Divide

85	24	63	45	17	31	96	50	
85	24	63	45	17	31	96	50	

Divide

85	5 24	63	45	17	31	96	50
85	24	63	45	17	31	96	50
85	24	63	45	17	31	96	5 50

Divide

	85	24	63	45	17	31	9	6	5	D	
8	5	24	63	45	17	31	-	96		50	
85	5	24	63	45	17	31	•	9	6	50)
85		24	63	45	17	32	1	9	6	5	0

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□ Sort & Merge



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□ Sort & Merge



15 **17 31 96 50**

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□ Sort & Merge

24	85						
85	24	63	45	17	31	96	50

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□ Sort & Merge



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□ Sort & Merge



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□ Sort & Merge



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□ Sort & Merge



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□ Sort & Merge

24	85	45	63	17	31		
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24			
----	--	--	--

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	45		
----	----	--	--

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	45	63	
----	----	----	--

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	45	63	85
----	----	----	----

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

24	45	63	85	17	31	50	96
24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Sort & Merge

17 24 31	45	50	63	85	96	
----------	----	----	----	----	----	--

24	45	63	85		17	31	50	96	
----	----	----	----	--	----	----	----	----	--

24	85	45	63	17	31	50	96
85	24	63	45	17	31	96	50

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□ Array is now sorted

17	24	31	45	50	63	85	96
----	----	----	----	----	----	----	----

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Merge Sort

Example 2



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Merge Sort

```
def merge(arr, 1, m, r):
   n1 = m - 1 + 1
   n2 = r - m
   L = [0] * (n1)
   R = [0] * (n2)
   for i in range(0 , n1):
       L[i] = arr[1 + i]
   for j in range(0 , n2):
       R[j] = arr[m + 1 + j]
   i = 0 # Initial index of first subarray
   j = 0 # Initial index of second subarray
   k = 1  # Initial index of merged subarray
   while i < n1 and j < n2 :
       if L[i] <= R[j]:
           arr[k] = L[i]
           i += 1
       else:
           arr[k] = R[j]
           j += 1
        k += 1
   while i < n1: # Copy the remaining elements of L[]
       arr[k] = L[i]
       i += 1
       k += 1
   while j < n2: # Copy the remaining elements of R[]
       arr[k] = R[j]
       j += 1
        k += 1
```

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```
def mergeSort(arr,l,r):
    if l < r:
        m = int((l+(r-1))/2)
        mergeSort(arr, l, m)
        mergeSort(arr, m+1, r)
        merge(arr, l, m, r)
    return arr</pre>
```

arr = [12, 11, 13, 5, 6, 7] print(mergeSort(arr,0,len(arr)-1))

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□Time Complexity: O(n * log(n))

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